## Quantum Processes and Computation Assignment 11, Wednesday, May 8, 2019

Exercise teachers:

Aleks Kissinger (aleks@cs.ru.nl) John van de Wetering (wetering@cs.ru.nl)

Handing in your answers: There are two options:

- 1. Deliver a hard copy to the mailbox of John van de Wetering. Mercator 1, 3rd floor.
- 2. E-mail a PDF to wetering@cs.ru.nl. Please include your name and the exercise number in the filename, e.g. ACHTERNAAM-qpc-exercise1.pdf.

Deadline: Tuesday, May 14, 12:00

**Goals:** After doing these exercises you know about a large class of quantum gates and you can reason with them. The total number of points is 100, distributed over 3 exercises. Material covered in book: sections 9.4, 12.1.

The Hadamard gate was defined as  $\begin{array}{c} \\ \\ \\ \end{array}$  := (

There are a few other equivalent definitions. Exercise 1 (9.112) (40 points):

(i) Follow the instructions of Exercise 9.112 in the book to show that:



There are a couple of sets quantum gates that crop up a lot in quantum computation. One of these

$$\mathbf{\dot{\varphi}} - \mathbf{\alpha} - \mathbf{\dot{\varphi}} = \begin{pmatrix} \mathbf{\dot{\alpha}} \\ \mathbf{\dot{\beta}} \\ \mathbf{\dot{\beta}} \\ \mathbf{\dot{\beta}} \end{pmatrix} \text{ and indeed } \mathbf{\dot{\varphi}} - \mathbf{\dot{\alpha}} - \mathbf{\dot{\varphi}} = \mathbf{\dot{\beta}} \quad \mathbf{\dot{\beta}} - \mathbf{\dot{\alpha}} - \mathbf{\dot{\beta}} = \mathbf{\dot{\beta}} \quad \mathbf{\dot{\beta}} - \mathbf{\dot{\beta}} + \mathbf{\dot{\beta}} = \mathbf{\dot{\beta}} \quad \mathbf{\dot{\beta}} - \mathbf{\dot{\beta}} + \mathbf{\dot{\beta}} = \mathbf{\dot{\beta}} \quad \mathbf{\dot{\beta}} - \mathbf{\dot{\beta}} + \mathbf{\dot{\beta}} = \mathbf{\dot{\beta}} \quad \mathbf{\dot{\beta}} = \mathbf{\dot{\beta}} = \mathbf{\dot{\beta}} \quad \mathbf{\dot{\beta}} = \mathbf{\dot{\beta}} = \mathbf{\dot{\beta}} \quad \mathbf{\dot{\beta}} = \mathbf{\dot{\beta}}$$

**Exercise 2 (12.8) (20 points):** Show using the ZX-calculus that the  $CZ(\alpha)$ -gate can be built from CNOTs and  $\circ$  phase gates as follows:



Recall that a single qubit unitary can be written as a succession of phase gates using a Euler Decomposition of the unitary. Since we have controlled phase gates in the ZX-calculus, we therefore have controlled single qubit unitaries. CNOT is a 2-qubit gate. Can we also make a controlled version of it? This would then be a 'controlled controlled not' gate, more commonly known as a *Toffoli*-gate. This is a three qubit gate that performs a NOT operation on the third qubit if the first two are in the  $|1\rangle$  state. This process of course involves an AND-gate. Constructing this gate is surprisingly involved.

**Exercise 3 (12.10) (40 points):** By evaluating on the Z-basis states (0, 1) and  $(\pi, \pi)$  show that the  $\wedge$  map below acts as a AND-gate.



You will need to use the following rules (that we haven't proven) and their colour reverses:



With the AND-gate in hand we can define the Toffoli-gate:



The Toffoli-gate is a very powerful gate. As is shown in Section 12.1.3 of the book, any unitary can be implemented in the ZX-calculus with the help of the Toffoli-gate. Somewhat surprisingly, it is the only gate you need for doing reversible classical computation, and approximate universal quantum computation can be achieved using just Toffoli- and Hadamard-gates.