Course: Type Theory and Coq

Exercises on Normalization

See the Slides and the Course Notes by Herman Geuvers for the definitions.

- 1. In the proof of WN for $\lambda \rightarrow$, the height of a type $h(\sigma)$ is defined by
 - $h(\alpha) := 0$
 - $h(\sigma_1 \rightarrow \ldots \rightarrow \sigma_n \rightarrow \alpha) := \max(h(\sigma_1), \ldots, h(\sigma_n)) + 1.$

Prove that this is the same as taking as the second clause

- $h(\sigma \rightarrow \tau) := \max(h(\sigma) + 1, h(\tau)).$
- 2. In the proof of WN for $\lambda \rightarrow$, it is stated that, if $M \longrightarrow_{\beta} N$ by contracting a redex of maximum height, h(M), that is not contained in another redex of maximum height, then this does not create a new redex of maximum height.

Show that this holds for the case

$$M = (\lambda x : A.x (\lambda v : B.x \mathbf{I}))(\lambda z : C.z (\mathbf{II}))$$

$$\longrightarrow_{\beta} (\lambda z : C.z (\mathbf{II}))(\lambda v : B.(\lambda z : C.z (\mathbf{II})) \mathbf{I}) = N$$

where $B = \alpha \rightarrow \alpha$, $C = B \rightarrow B$ and $A = C \rightarrow B$.

Also show that $m(M) >_l m(N)$.

- 3. Prove that type reduction is SN for $\lambda 2$ a la Church. (Define a simple measure on terms that decreases with type reduction.)
- 4. Prove for that for $A, B \in SAT$, $A \to B \in SAT$. (Here, $A \to B := \{M | \forall N \in A(M \ N \in B)\}$. Check the slides or course notes for the definition of SAT, the collection of saturated sets.)
- 5. The Soundness of the saturated sets model for $\lambda 2$ is proved by induction on the derivation. Do the case for the \forall -introduction rule.