

# statistics on digital libraries of mathematics

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<http://www.cs.ru.nl/~freek/notes/stats.pdf>

<http://www.cs.ru.nl/~freek/talks/stats.pdf>

## libraries of mathematics

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- **text**

- paper  
journals, textbooks
- electrons  
archives, websites

- **code**

- mathematical tables
- software  
numerical, computer algebra, subject specific
- formalizations

## 80 out of 100 theorems









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1. The Irrationality of the Square Root of 2  $\geq 17$
2. Fundamental Theorem of Algebra 4
3. The Denumerability of the Rational Numbers 6
4. Pythagorean Theorem 6
5. Prime Number Theorem 2
6. Gödel's Incompleteness Theorem 3
7. Law of Quadratic Reciprocity 4
8. The Impossibility of Trisecting the Angle and Doubling the Cube 1
9. The Area of a Circle 1
10. Euler's Generalization of Fermat's Little Theorem 4
11. The Infinitude of Primes 6
12. The Independence of the Parallel Postulate 0
13. Polyhedron Formula 1
- ...

## proof assistants

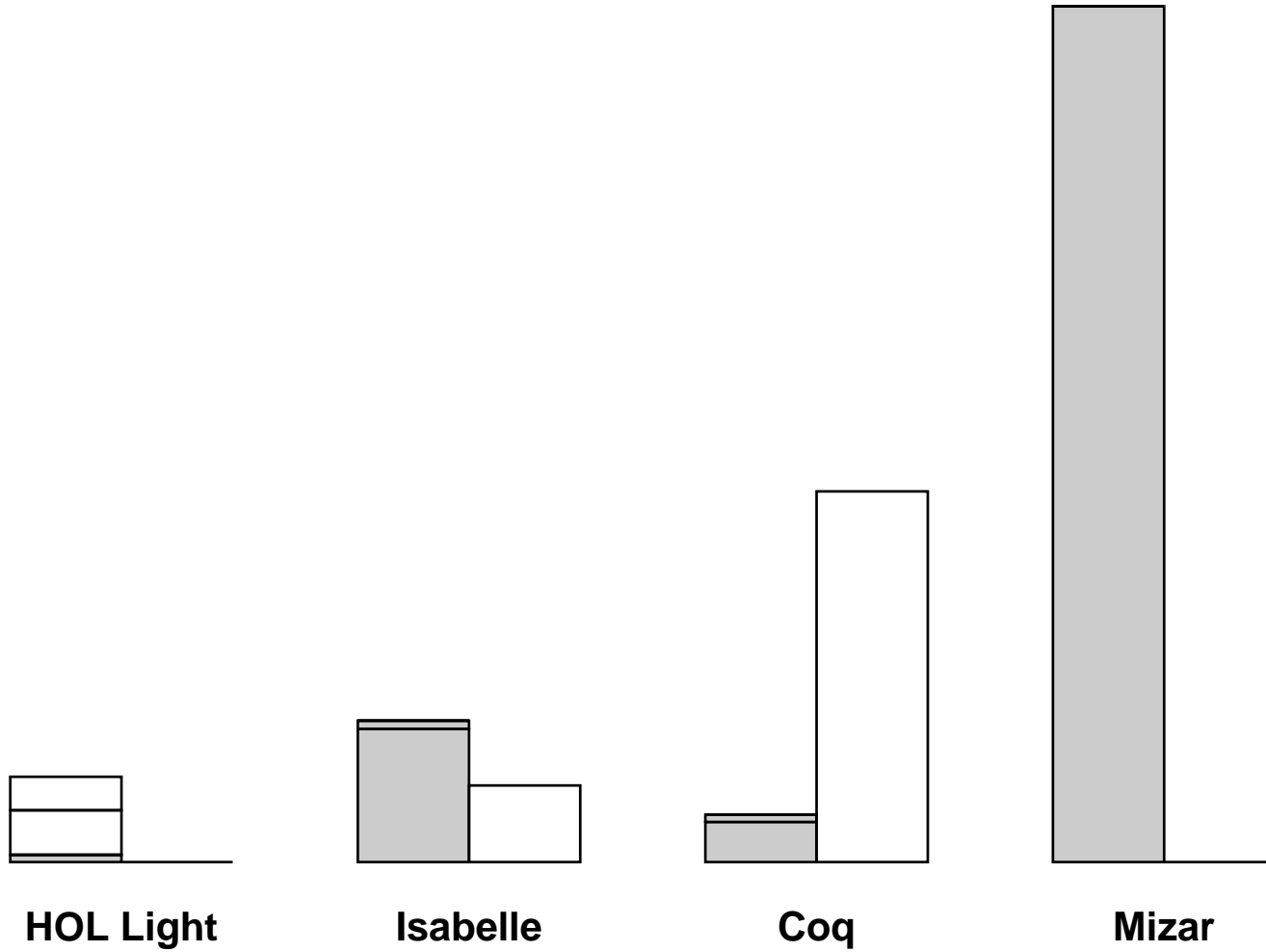
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only **five systems** seriously used **for mathematics**

HOL {	 	<b>HOL Light</b>	69
		<b>ProofPower</b>	42
	 	<b>Isabelle</b>	40
		<b>Coq</b>	39
	 	<b>Mizar</b>	45

## library sizes

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## content of the files

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formalization  $\approx$  long chain of **lemmas**

lemma  $\approx$

label +

statement +

**proof**

```

C (* ===== *)
C (* Binary expansions as a bijection between numbers and finite sets. *)
C (* ===== *)
B
T let LT_POW2_REFL = prove
T ('!n. n < 2 EXP n',
P  INDUCT_TAC THEN REWRITE_TAC[EXP] THEN TRY(POP_ASSUM MP_TAC) THEN ARITH_TAC);;
B
T let BINARY_INDUCT = prove
T ('!P. P 0 /\ (!n. P n ==> P(2 * n) /\ P(2 * n + 1)) ==> !n. P n',
P  GEN_TAC THEN STRIP_TAC THEN MATCH_MP_TAC num_WF THEN GEN_TAC THEN
P  STRIP_ASSUME_TAC(ARITH_RULE
P   'n = 0 \/ n DIV 2 < n /\ (n = 2 * n DIV 2 \/ n = 2 * n DIV 2 + 1)') THEN
P  ASM_MESON_TAC[]);;
B
T let BOUNDED_FINITE = prove
T ('!s. (!x:num. x IN s ==> x <= n) ==> FINITE s',
P  REPEAT STRIP_TAC THEN MATCH_MP_TAC FINITE_SUBSET THEN EXISTS_TAC '0..n' THEN
P  ASM_SIMP_TAC[SUBSET; IN_NUMSEG; FINITE_NUMSEG; LE_0]);;
B
T let EVEN_NSUM = prove
T ('!s. FINITE s /\ (!i. i IN s ==> EVEN(f i)) ==> EVEN(nsum s f)',
P  REWRITE_TAC[GSYM IMP_IMP] THEN MATCH_MP_TAC FINITE_INDUCT_STRONG THEN

```

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```

C (* Title:      HOL/Quadratic_Reciprocity/Gauss.thy
C   ID:          $Id: Int2.thy,v 1.12 2007/06/11 09:06:07 chaieb Exp $
C   Authors:     Jeremy Avigad, David Gray, and Adam Kramer
C *)
B
E header {*Integers: Divisibility and Congruences*}
B
S theory Int2 imports Finite2 WilsonRuss begin
B
D definition
D   MultInv :: "int => int => int" where
D   "MultInv p x = x ^ nat (p - 2)"
B
B
E subsection {* Useful lemmas about dvd and powers *}
B
T lemma zpower_zdvd_prop1:
T   "0 < n \<Longrightarrow> p dvd y \<Longrightarrow> p dvd ((y::int) ^ n)"
P   by (induct n) (auto simp add: zdvd_zmult zdvd_zmult2 [of p y])
B
T lemma zdvd_bounds: "n dvd m ==> m \<le> (0::int) | n \<le> m"
P proof -
P   assume "n dvd m"

```



```

C (*****
C (* v      * The Coq Proof Assistant / The Coq Development Team *)
C (* <O___,, * CNRS-Ecole Polytechnique-INRIA Futurs-Universite Paris Sud *)
C (* \VV/ *****
C (* // * This file is distributed under the terms of the *)
C (* * GNU Lesser General Public License Version 2.1 *)
C (*****
C (*i $Id: Decidable.v 5920 2004-07-16 20:01:26Z herbelin $ i*)
B
E (** Properties of decidable propositions *)
B
D Definition decidable (P:Prop) := P \/ ~ P.
B
T Theorem dec_not_not : forall P:Prop, decidable P -> (~ P -> False) -> P.
P unfold decidable in |- *; tauto.
P Qed.
B
T Theorem dec_True : decidable True.
P unfold decidable in |- *; auto.
P Qed.
B
T Theorem dec_False : decidable False.
P unfold decidable, not in |- *; auto.

```

```
C :: Non negative real numbers. Part II
C ::  by Andrzej Trybulec
C ::
C :: Received March 7, 1998
C :: Copyright (c) 1998 Association of Mizar Users
B
S environ
B
S vocabularies ARYTM_2, BOOLE, ORDINAL2, ARYTM_3, ARYTM_1;
S notations TARSKI, SUBSET_1, ARYTM_3, ARYTM_2;
S constructors ARYTM_2;
S requirements SUBSET;
S theorems ARYTM_2;
B
S begin
B
L reserve x,y,z for Element of REAL+;
B
T theorem Th1:
T    $x + y = y$  implies  $x = \{\}$ 
P proof reconsider o =  $\{\}$  as Element of REAL+ by ARYTM_2:21;
P   assume  $x + y = y$ ;
P   then  $x + y = y + o$  by ARYTM_2:def 8;
```

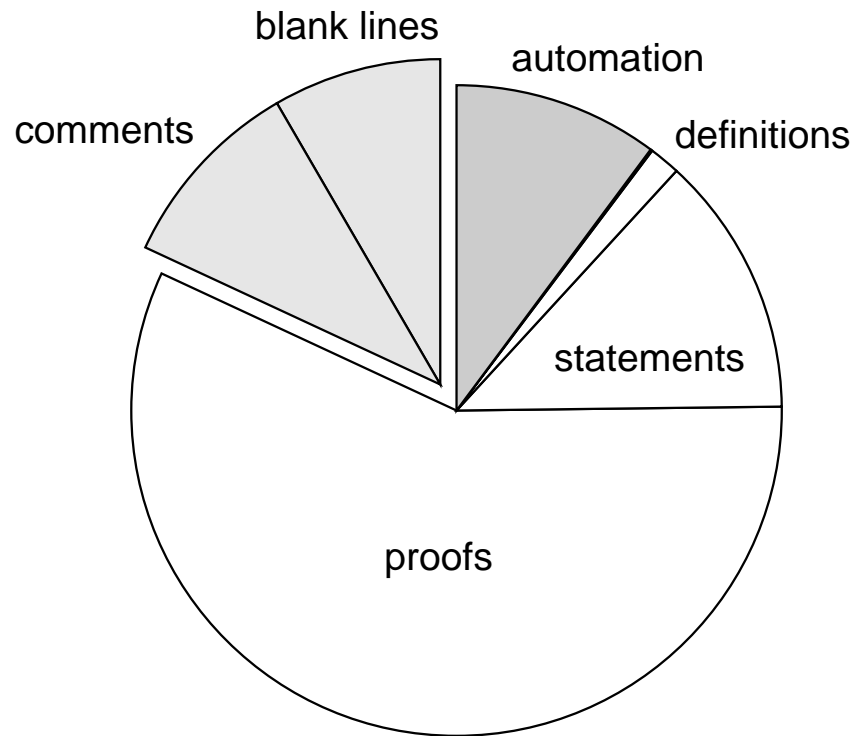
## 11 kinds of lines

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- B blank lines
- C comments
- E documentation
- S modules: imports and sectioning
- L contexts
- D definitions
- N notation
- H automation: directives
- X automation: program code
- T theorem statements
- P proofs

## relative sizes: HOL Light

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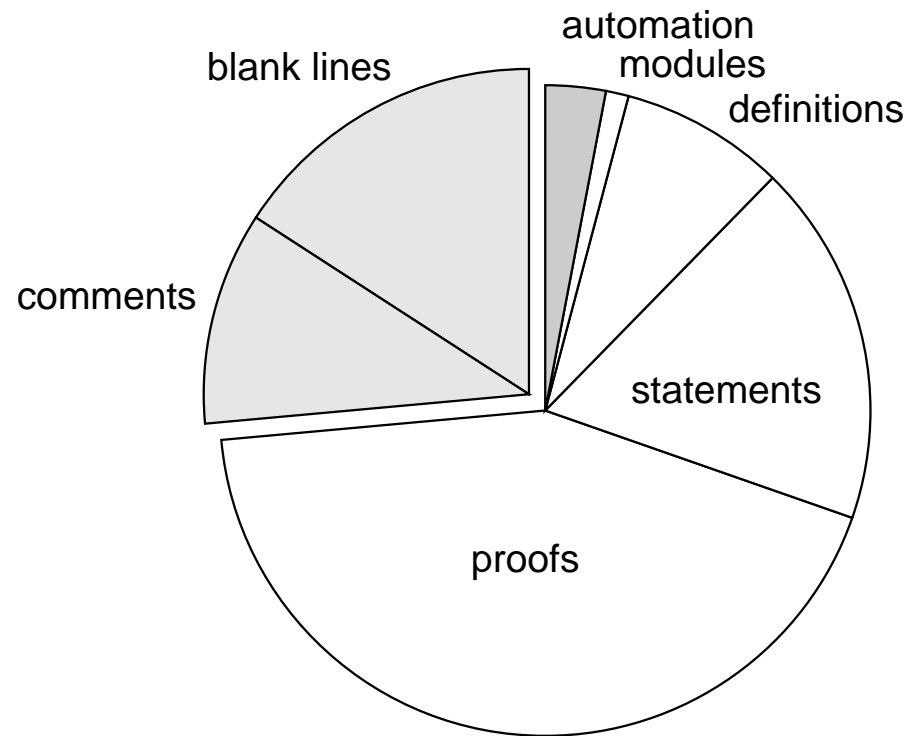


almost no module lines

**small definitions part!**

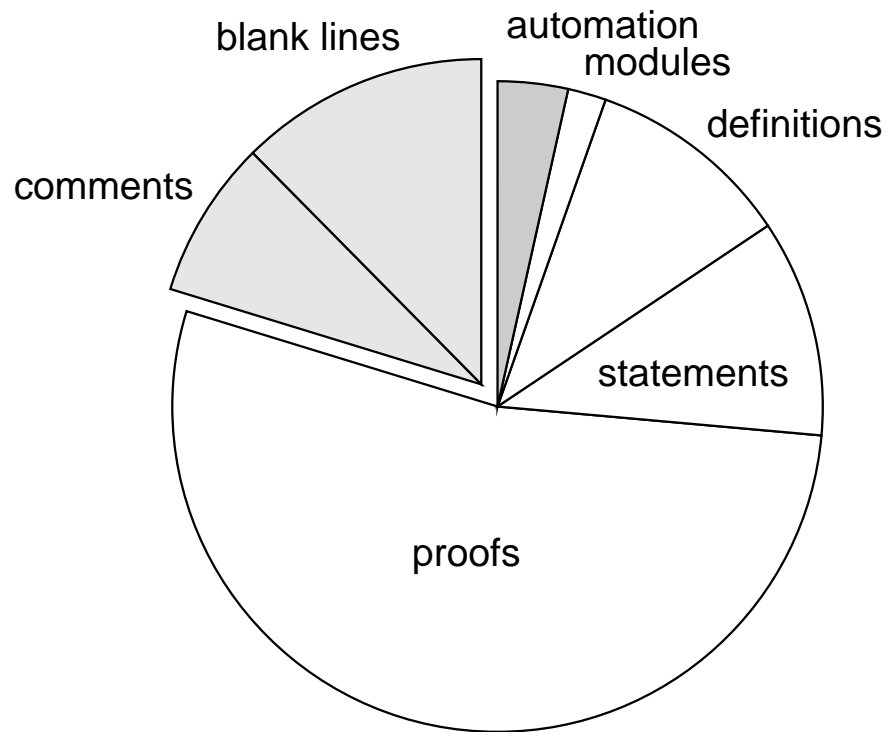
## relative sizes: Isabelle

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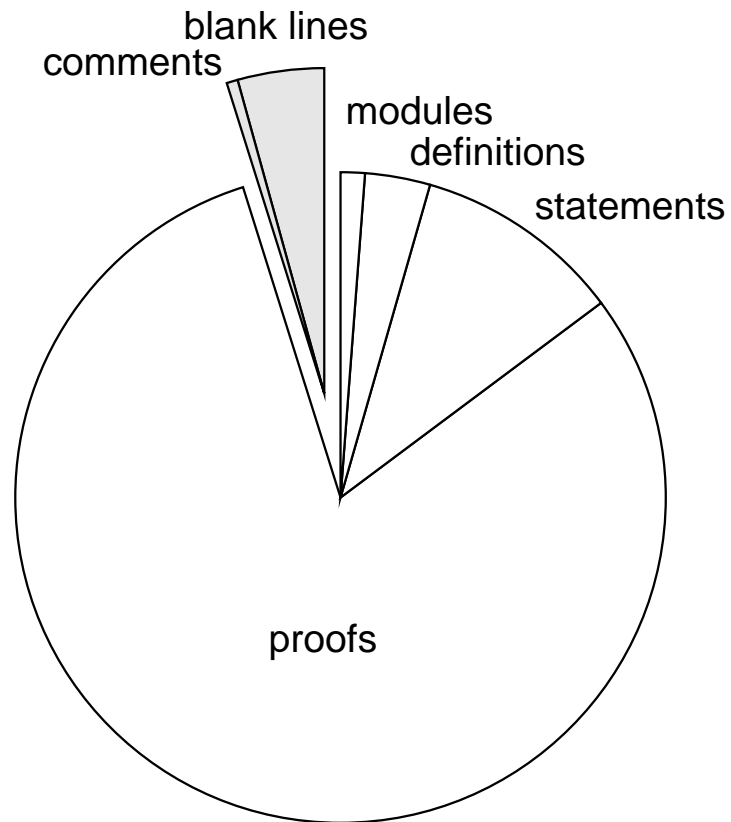
## relative sizes: Coq

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## relative sizes: Mizar

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no automation lines  
almost no comments  
**large proofs part!**

## lessons learned

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- formalizations all are very similar
  - ... despite fundamental differences between proof assistants (foundations, interaction styles)
- formalizations consist primarily of proofs
- classification of line types in formalizations
- small definitions are good!
- proof automation is important!