

1 Counting Letters

- a) Let A be a non-empty, finite alphabet and $a \in A$ a letter.

Define by structural induction a map

(1pt)

$$|\cdot|_a: A^* \rightarrow \mathbb{N}$$

that counts the number of occurrences of the letter a in a word.

Begin Secret Info:

Solution

We define $|w|_a$ by induction on w :

- Base case: $|\lambda|_a = 0$.
- Step case for $xv \in A^*$. So we have already defined $|v|_a$. We distinguish two cases: If $x = a$, then $|av|_a = |v|_a + 1$, otherwise if $x \neq a$, then we define $|xv|_a = |v|_a$.

End Secret Info

- b) Show, by induction, that for any two words $w, u \in A^*$

(2pt)

$$|wu|_a = |w|_a + |u|_a.$$

Begin Secret Info:

Solution

We show $|wu|_a = |w|_a + |u|_a$ by induction on w .

- Base case: $w = \lambda$. Then we get

$$|w|_a = |u|_a = 0 + |u|_a = |w|_a + |u|_a.$$

- Induction step: Let $w = xv$ for some $x \in A$. Induction Hypothesis: $|vu|_a = |v|_a + |u|_a$ holds for all $u \in A^*$. We distinguish cases:

If $x = a$, then

$$|avu|_a = 1 + |vu|_a \stackrel{IH}{=} 1 + |v|_a + |u|_a = |av|_a + |u|_a = |w|_a + |u|_a.$$

If $x \neq a$, then

$$|xvu|_a = |vu|_a \stackrel{IH}{=} |v|_a + |u|_a = |xv|_a + |u|_a = |w|_a + |u|_a.$$

This induction shows, that $|wu|_a = |w|_a + |u|_a$ for all $w, u \in A^*$.

End Secret Info

2 Regular expression

Let L be the language given by

$$\{w \in \{a, b\}^* \mid \text{every } b \text{ in } w \text{ is directly followed by an } a\}$$

Give a regular expression for the language L and explain your answer.

(1pt)

Begin Secret Info:

Solution

We use $a^*(baa^*)^*$.

Explanation: A word in the language should consist of a finite (possibly zero!) number of ba 's that can be interleaved with a 's, so a word in L is of the shape $a^{n_1}(ba)a^{n_2} \dots (ba)a^{n_k}$ with $n_1, n_2 \dots n_k \geq 0$. This is captured by the regular expression $a^*(baa^*)^*$.

Obvious alternatives: $a^*(a^*baa^*)^*$ or $(a^*ba)^*a^*$.

End Secret Info