

## Exercises Coalgebra for Lecture 5

The exercises labeled with (\*) are optional and more advanced.

- Let  $A$  be a set, and consider the functor  $F: \mathbf{Set} \rightarrow \mathbf{Set}$ , defined on a set  $X$  by  $F(X) = A \times X \times X$  and on a function  $f$  by  $F(f) = \text{id}_A \times f \times f$ . An *infinite binary tree* (node-labelled in  $A$ ) is a function  $t: \{l, r\}^* \rightarrow A$ , where  $\{l, r\}^*$  is the set of words over  $l, r$ . The empty word is denoted by  $\varepsilon \in \{l, r\}^*$ . The set  $A^{\{l, r\}^*}$  of all infinite binary trees is denoted by  $T$ .

Given a tree  $t \in T$ , we define  $t_l, t_r \in T$  as follows:  $t_l(w) = t(lw)$  and  $t_r(w) = t(rw)$ , for all  $w \in \{l, r\}^*$ . Consider the  $F$ -coalgebra  $z: T \rightarrow A \times T \times T$  defined by

$$z(t) = (t(\varepsilon), t_l, t_r)$$

The aim of this exercise is to show that  $(T, z)$  is a final  $F$ -coalgebra.

- Describe, in words, what  $t_l$  and  $t_r$  are, given a tree  $t \in T$ .
  - Let  $\langle g_\varepsilon, g_1, g_2 \rangle: X \rightarrow A \times X \times X$  be an  $F$ -coalgebra. Show that a map  $h: X \rightarrow T$  is a homomorphism from  $(X, \langle g_\varepsilon, g_1, g_2 \rangle)$  to  $(T, z)$  if and only if for all  $x \in X$ :
    - $h(x)(\varepsilon) = g_\varepsilon(x)$
    - $h(x)_l = h(g_1(x))$
    - $h(x)_r = h(g_2(x))$
  - Conclude that such a homomorphism (as in the previous exercise) exists (why?).
  - Show that if  $h, k$  are both homomorphisms from  $(X, g)$  to  $(T, z)$  then  $h = k$ , by proving by induction on  $w \in \{l, r\}^*$  that for all  $x \in X$ :  $h(x)(w) = k(x)(w)$ .
- Consider the functor  $F: \mathbf{Set} \rightarrow \mathbf{Set}$ , defined on a set  $X$  by  $F(X) = X + A$ , where  $A$  is fixed set.
    - Make an (educated) guess about a final coalgebra for  $F$ . How would you define the unique homomorphism  $\text{beh}$ ?
    - (\*) Prove that your answer to the previous question is correct.
  - Let  $F: \mathcal{C} \rightarrow \mathcal{C}$  be an endofunctor on a category  $\mathcal{C}$ . In the lecture, we defined the category  $\mathbf{CoAlg}(F)$  whose objects are  $F$ -coalgebras, and whose morphisms are coalgebra morphisms. Show that this is indeed a category, by checking the necessary axioms.
  - In the lecture, we defined, for any given automaton  $(S, \langle \epsilon, \delta \rangle)$ , a map  $\text{beh}: S \rightarrow 2^{A^*}$ .  
Finish the proof that  $(2^{A^*}, \langle e, d \rangle)$  is a final coalgebra, by showing that  $\text{beh}$  is the unique coalgebra homomorphism from  $(S, \langle \epsilon, \delta \rangle)$  to  $(2^{A^*}, \langle e, d \rangle)$ .

5. (\*) Let  $F: \mathbf{Set} \rightarrow \mathbf{Set}$  be a functor with a final coalgebra  $(Z, z)$ . In the lecture, we defined two states  $x, y \in X$  of an  $F$ -coalgebra  $(X, f)$  to be *behaviourally equivalent* if  $\text{beh}(x) = \text{beh}(y)$ , where  $\text{beh}$  is the unique homomorphism from  $(X, f)$  to  $(Z, z)$ . Show that  $\text{beh}(x) = \text{beh}(y)$  if and only if there exists an  $F$ -coalgebra  $(Y, g)$  and a homomorphism  $h: X \rightarrow Y$  from  $(X, f)$  to  $(Y, g)$  such that  $h(x) = h(y)$ . Hint: draw a suitable diagram to clarify the situation.
6. (\*) We would like to define a functor  $S: \mathbf{Set} \rightarrow \mathbf{Set}$  by  $S(X) = X^\omega$ , i.e., a set  $X$  is mapped to the set of streams over  $X$ .
- (a) Define  $S$  on a function  $f: X \rightarrow Y$ , using that  $Y^\omega$  is the final stream system over  $Y$ ;  $S(f)$  should apply  $f$  to all elements of a given stream.
  - (b) Show that  $S$  is functorial.
7. (\*) Let  $F: \mathcal{C} \rightarrow \mathcal{C}$  be a functor on a category  $\mathcal{C}$ . Suppose  $\mathcal{C}$  has an initial object, and a coproduct  $X + Y$  for any objects  $X, Y$ . Show that  $\text{CoAlg}(F)$  has an initial object and all coproducts as well.