## **Consistency-based diagnosis (cont.)**



Difference between predicted behaviour and observed behaviour  $\Rightarrow$  defect!

#### Originators:

- R. Reiter, "A Theory of diagnosis from first principles", Artificial Intelligence, vol. 32, 57–95, 1987.
- J. de Kleer, A.K. Macworth, and R. Reiter, "Characterising diagnoses and systems", Artificial Intelligence, vol. 52, 197–222, 1992.

### **Plan for today**

- Quick revision of basic concepts
- There are some opimisations possible for the hitting-set algorithm
- Consistency-based diagnosis is an example of non-monotonic reasoning (quite common in AI). We show why this is the case
- Finally, dagnoses can be seen as hypotheses that are revised when new observations are made. We extend the theory in this way

 $\implies$  First, revision!

## **System specification**



Multiplier-adder

#### **SYStem specification** SYS = (SD, COMPS):

SD (System Description):

 $\forall x ((\mathsf{MUL}(x) \land \neg \mathsf{Ab}(x)) \to \operatorname{in}_1(x) \times \operatorname{in}_2(x) = \operatorname{out}(x)) \\ \forall x ((\mathsf{ADD}(x) \land \neg \mathsf{Ab}(x)) \to \operatorname{in}_1(x) + \operatorname{in}_2(x) = \operatorname{out}(x)) \\ \operatorname{MUL}(M_1), \operatorname{MUL}(M_2), \operatorname{MUL}(M_3), \operatorname{ADD}(A_1), \operatorname{ADD}(A_2) \\ \operatorname{in}_1(A_1) = \operatorname{out}(M_1), \operatorname{in}_2(A_1) = \operatorname{out}(M_2) \\ \operatorname{in}_1(A_2) = \operatorname{out}(M_2), \operatorname{in}_2(A_2) = \operatorname{out}(M_3)$ 

• COMPS = { $M_1, M_2, M_3, A_1, A_2$ }

## **Diagnostic problem**



- **System specification** SYS = (SD, COMPS)
- Diagnostic problem DP = (SYS, OBS), with OBS a set of observations

OBS = {
$$in_1(M_1) = 3$$
,  $in_2(M_1) = 2$ ,  $in_1(M_2) = 3$ ,  $in_2(M_2) = 2$ ,  
 $in_1(M_3) = 2$ ,  $in_2(M_3) = 3$ ,  $out(A_1) = 10$ ,  $out(A_2) = 12$ }

# Diagnosis

- Diagnostic problem DP = (SYS, OBS), with OBS a set of observations
- A diagnosis D: smallest (subset minimal) set of components, such that

Faulty

components

 $SD \cup OBS \cup \{Ab(c) \mid c \in D\} \cup \{\neg Ab(c) \mid c \in COMPS - D\}$ 

Nonfaulty

components

is consistent

• Remark:  ${Ab(c) | c \in D}$  can be omitted (why?)

For the multiplier-adder:  $D = \{A_1\}, \{M_1\}, \{M_2, M_3\}, \{A_2, M_2\}$ 

### **Conflict set and hitting set**

Let  $\mathrm{CS}\subseteq\mathrm{COMPS}$  be a set of components, then  $\overline{\mathrm{CS}}$  is called a conflict set iff

$$SD \cup OBS \cup \{\neg Ab(c) \mid c \in CS\}$$

is inconsistent

**Proposition:** For each  $D \subseteq \text{COMPS}$  that is a diagnosis and each conflict set CS it holds that:  $D \cap \text{CS} \neq \emptyset$ 

Theorem: *D* is a diagnosis for diagnostic problem DP = (SYS, OBS) iff *D* is a minimal hitting set for all conflict sets of DP

### **Hitting-set tree**

Let *F* be a set of sets

• Let  $T = (V, E, l_V, l_E)$  be a labelled tree, with

node labels

$$l_V(v) = \begin{cases} S & \text{if } S \in F, \, S \neq \emptyset \\ \checkmark & \text{otherwise} \end{cases}$$

and

• edge labels if  $l_V(v) = S$  then  $\forall s \in S$ :  $l_E(v, v_s) = s$ 

### **Hitting sets**

The hitting set H(v) for node v is defined as:

 $H(v) = \{l_E(u, w) \mid (u, w) \text{ is on the path from the root to } v\}$ 

Example (incorrect why?):

 $\begin{array}{c|c} a \{1,3,5\} \\ 1 & 3 & 5 \\ \hline b \{2\} & c \{4\} & \checkmark \\ c & 4 & \checkmark \\ c & f & \checkmark \\ c & \land \\$ 

$$l_V(a) = \{1, 3, 5\},$$
  
 $l_V(b) = \{2\}, l_V(e) = \checkmark,$  etc.
 $l_E(a, b) = 1, l_E(a, c) = 3,$   
 $l_E(b, e) = 2,$  etc.

• 
$$H(a) = \emptyset$$

• 
$$H(b) = \{1\}$$

• 
$$H(e) = \{1, 2\}$$

•  $H(f) = \{3, 4\}$ 

### **Hitting sets**

The hitting set H(v) for node v is defined as:

 $H(v) = \{l_E(u, w) \mid (u, w) \text{ is on the path from the root to } v\}$ Example (correct):

 $l_V(a) = \{1, 3, 5\},$  $a \{1, 3, 5\}$  $l_V(b) = \{2\}, l_V(e) = \checkmark,$  etc.  $|\mathbf{1/3}|$  5  $\mathbf{I}_E(a,b) = 1, \ l_E(a,c) = 3,$  $b\left\{2\right\} c\left\{4\right\} d$  $l_{E}(b,e) = 2$ , etc.  $I(a) = \emptyset$ e $= H(b) = \{1\}$  $H(e) = \{1, 2\}$ g $H(f) = \{3, 4\}$ 

#### **Example: full-adder**



### **Example HS tree**



1.  $CS_1$ TP(SD, COMPS, OBS);  $CS_1 \leftarrow \{X_1, X_2\}$ 

2.  $CS_2 \leftarrow TP(SD, COMPS - \{X_1\}, OBS);$  $CS_2 \leftarrow \checkmark$  (diagnosis found)

**3.**  $CS_3 \leftarrow TP(SD, COMPS - \{X_2\}, OBS);$  $CS_3 \leftarrow \{X_1, A_2, O_1\}$ 

4. :

Diagnoses *D*:  $\{X_1\}$ ,  $\{X_2, A_2\}$ ,  $\{X_2, O_1\}$ (note that  $\{X_2, X_1\}$  not subset minimal)

### **Pruning of the hitting-set tree**



Note that there is no need to extend the hitting set  $H(v_4)$  (as  $\{X_2, X_1\}$  is not subset minimal)

 $\Rightarrow$  pruning of the hitting-set tree

### **Hitting-set tree**



Minimal hitting set: H(v) is subset-minimal and  $l_V(v) = \checkmark$ 

**Examples:**  $H(v_6) = \{1, 4\} \subset H(v_{11}); H(v_7) = \{2, 4\} \subset H(v_{14})$ 

### **Optimisation**

Pruning the HS tree:

- 1. H(v) = H(v'): prune the subtree with root v' and  $l_V(v') = \times$
- 2.  $H(v) \subset H(v')$ : ignore v',  $l_V(v') = \times$
- 3. for  $l_V(v) = S$ ,  $l_V(v') = S' \in F$  with  $S' \subset S$ :



• Reuse of labels when F is dynamic (as in diagnosis): if  $S' \in F$  and  $H(v) \cap S' = \emptyset$ , then  $l_V(v) = S'$ 









### **Non-monotonic reasoning**

- **Knowledge base** KB
- Add knowledge to KB and obtain new knowledge base KB'
- If KB ⊢ Results and KB' ⊢ Results' then Results ⊆ Results' does not hold in general ⇒ more knowledge does not always yield more results
- Consistency-based reasoning is an example of non-monotonic reasoning. Why?

 $SD \cup OBS \cup \{\neg Ab(c) \mid c \in COMPS - D\} \nvDash \bot$ 

e.g., larger OBS or SD may make D smaller or different

## **Default logic**

DT = (W, R) is a default theory, where W = {Elephant(john)}, i.e., John is an elephant, and the following default R:

$$\frac{\text{Elephant}(x):\text{Grey}(x)}{\text{Grey}(x)}$$

If being grey is consistent with our knowledge, conclude 'grey', so conclude Grey(john)



general form default

prerequisite : justifications consequent

## **Reasoning in default logic**

Let DT = (W, R) be a default theory (W a set of logical formulas and a set of defaults R):

- E = Th(E) (so-called fixed point)
- E includes the maximal set of conclusions obtained by applying defaults in R

If 
$$\frac{A:B_1,\ldots,B_n}{C} \in R$$
,  $A \in E$  and  $\neg B_1,\ldots,\neg B_n \notin W$ , then  $C \in E$ 

E is called an extension and Th is the derivation operator (deduction + default rule application)

DT = (W, R), where

```
W = \{ Elephant(clyde), \neg Grey(john) \}
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i.e., Clyde is an elephant and John is not grey, and the following default R:

$$\frac{\text{Elephant}(x):\text{Grey}(x)}{\text{Grey}(x)}$$

'elephants are normally grey'

**Extension:**  $E = \{ Elephant(clyde), Grey(clyde), \neg Grey(john) \}$ 

## **Diagnosis as non-monotonic reasoning**

- Map a diagnostic problem to a default theory DT
- A diagnosis D predicts a formula  $\varphi$  iff

 $\mathrm{SD} \cup \mathrm{OBS} \cup \{\neg \mathrm{Ab}(c) \mid c \in \mathrm{COMPS} - D\} \cup \{\mathrm{Ab}(c) \mid c \in D\} \vDash \varphi$ 

• Lemma: DT = (W, R) is a default theory with extension E iff

$$E = \operatorname{Th} \left( W \cup \{ L \mid : L/L \in \Delta \} \right)$$

with subset-maximal set of defaults  $\Delta \subseteq R$  such that

$$W \cup \{L \mid : L/L \in \Delta\} \nvDash \bot$$

Remark:  $:\psi/\psi$  is a so-called normal default (default without prerequisite and justification that is the same as the conclusion)

## Logical characterisation of diagnosis

Theorem. Let  $\mathrm{DP}=(\mathrm{SYS},\mathrm{OBS})$  be a diagnostic problem. Let

$$DT = \left(SD \cup OBS, \left\{\frac{: \neg Ab(c)}{\neg Ab(c)} \mid c \in COMPS\right\}\right)$$

be a default theory with extension *E*, then *D* is a diagnosis for DP iff  $E = \{ \varphi \mid D \text{ predicts } \varphi \}$ 

Proof: *E* is an extension of DT, thus (Lemma):

$$\mathrm{SD} \cup \mathrm{OBS} \cup \left\{ \neg \mathrm{Ab}(c) \mid \frac{: \neg \mathrm{Ab}(c)}{\neg \mathrm{Ab}(c)} \in \Delta \right\} \nvDash \bot$$

with  $\Delta \subseteq R$  such that  $\Delta$  subset-maximal. Suppose that  $D = \{c \mid c \in \text{COMPS}, \frac{:\neg Ab(c)}{\neg Ab(c)} \notin \Delta\}$ , then . . .

## Logical characterisation of diagnosis

Theorem. Let  $\mathrm{DP}=(\mathrm{SYS},\mathrm{OBS})$  be a diagnostic problem. Let

$$\mathrm{DT} = \left(\mathrm{SD} \cup \mathrm{OBS}, \left\{\frac{: \neg \mathrm{Ab}(c)}{\neg \mathrm{Ab}(c)} \mid c \in \mathrm{COMPS}\right\}\right)$$

be a default theory with extension *E*, then *D* is a diagnosis for DP iff  $E = \{ \varphi \mid D \text{ predicts } \varphi \}$ 

Proof (continued):

$$\left\{\neg \operatorname{Ab}(c) \mid \frac{:\neg \operatorname{Ab}(c)}{\neg \operatorname{Ab}(c)} \in \Delta\right\} = \left\{\neg \operatorname{Ab}(c) \mid c \in \operatorname{COMPS} - D\right\}$$

Thus,  $E = \text{Th}(\text{SD} \cup \text{OBS} \cup \{\neg \text{Ab}(c) \mid c \in \text{COMPS} - D\})$ , and  $E = \{\varphi \mid D \text{ predicts } \varphi\}$ 

$$DP = (SYS, OBS), with$$
  

$$SD = \{\forall x ((ANDG(x) \land \neg Ab(x)) \rightarrow (out(x) = and(in_1(x), in_2(x)))), \\ \forall x ((XORG(x) \land \neg Ab(x)) \rightarrow (out(x) = xorg(in_1(x), in_2(x)))), \\ ANDG(A), XORG(X), \\ out(A) = in_1(X)\}$$

 $COMPS = \{A, X\} \text{ and } OBS = \{in_1(A) = 1, in_2(A) = 1, in_2(X) = 0, out(A) = 0, out(X) = 1\}$ 

Default theory  $DT = (SD \cup OBS, R)$ , with

$$R = \left\{ \frac{: \neg \operatorname{Ab}(A)}{\neg \operatorname{Ab}(A)}, \frac{: \neg \operatorname{Ab}(X)}{\neg \operatorname{Ab}(X)} \right\}$$

 $E = \text{Th}(\text{SD} \cup \text{OBS} \cup \{\neg \text{Ab}(X)\})$ , e.g.,  $\text{Ab}(A) \in E$  (we can predict that A is abnormal)

#### **Extra measurements**

Recall: a diagnosis D predicts a formula  $\varphi$  iff

 $\mathrm{SD} \cup \mathrm{OBS} \cup \{\neg \mathrm{Ab}(c) \mid c \in \mathrm{COMPS} - D\} \vDash \varphi$ 

if  $\varphi$  is a set of extra observations (measurements), then:

- 1. Every diagnosis *D* for DP = (SYS, OBS) that predicts  $\varphi$  is also a diagnosis for  $DP = (SYS, OBS \cup \{\varphi\})$ , i.e., the measurement  $\varphi$  confirms *D*
- 2. No diagnosis for DP = (SYS, OBS) that predicts  $\neg \varphi$  is also a diagnosis for  $DP = (SYS, OBS \cup \{\varphi\})$ , i.e., the measurment  $\varphi$  disconfirms D
- 3. Any diagnosis *D* for  $DP = (SYS, OBS \cup \{\varphi\})$  which is not a diagnosis for DP' = (SYS, OBS) is a strict superset of a diagnosis of DP' which predicts  $\neg \varphi$



- **Diagnostic problem** DP = (SYS, OBS), with
- Set of observations

OBS = {
$$in_1(M_1) = 3$$
,  $in_2(M_1) = 2$ ,  $in_1(M_2) = 3$ ,  $in_2(M_2) = 2$ ,  
 $in_1(M_3) = 2$ ,  $in_2(M_3) = 3$ ,  $out(A_1) = 10$ ,  $out(A_2) = 12$ }

• Predictions w.r.t.  $out(M_2)$ : diagnosis  $D_1 = \{M_1\}$  predicts  $out(M_2) = 6$ ,  $D_2 = \{A_1\}$  predicts  $out(M_2) = 6$ ,  $D_3 = \{M_2, M_3\}$  predicts  $out(M_2) = 4$ ,  $D_4 = \{M_2, A_2\}$ predicts  $out(M_2) = 4$ 

## **Example (continued)**



- **Diagnostic problem** DP = (SYS, OBS), with
- Set of observations (with new one on  $M_2$ ):

OBS = {
$$in_1(M_1) = 3$$
,  $in_2(M_1) = 2$ ,  $in_1(M_2) = 3$ ,  $in_2(M_2) = 2$ ,  
 $in_1(M_3) = 2$ ,  $in_2(M_3) = 3$ ,  $out(A_1) = 10$ ,  $out(A_2) = 12$ ,  
 $out(M_2) = 5$ }

The new observation  $out(M_2) = 5$  disconfirms all previous diagnoses

### **Example (continued)**



- **Diagnostic problem** DP = (SYS, OBS), with
- Set of observations (with new one on  $M_2$ ):

$$OBS = \{ in_1(M_1) = 3, in_2(M_1) = 2, in_1(M_2) = 3, in_2(M_2) = 2, in_1(M_3) = 2, in_2(M_3) = 3, out(A_1) = 10, out(A_2) = 12, out(M_2) = 5 \}$$

• New diagnoses:  $D'_1 = \{M_1, M_2, M_3\}, D'_2 = \{M_1, M_2, A_2\}, D'_3 = \{M_2, M_3, A_1\}, D'_4 = \{M_2, A_1, A_2\};$  note supersets of the old diagnoses, e.g.,  $D_1, D_3 \subseteq D'_1$  (case 3)

### **Definition of diagnosis revisited**

Example two inverters  $I_k$ , k = 1, 2:



• OBS =  $\{ in(I_1) = 0, out(I) = 1 \}$ 

• Diagnoses:  $\{I_1\}$  and  $\{I_2\}$ 

However, {I<sub>1</sub>, I<sub>2</sub>} might also be a diagnosis (only excluded because of definition using subset-minimality condition)

#### **Faul models**



 $SD = \{ \forall x ((INV(x) \land \neg Ab(x)) \rightarrow \neg (out(x) = in(x))), \\ \forall x ((INV(x) \land Ab(x)) \rightarrow (Stuck\_at0(x) \lor (out(x) = in(x)))), \\ \forall x (Stuck\_at0(x) \rightarrow out(x) = 0), \\ INV(I_1), \\ INV(I_2) \}$ 

Knowledge about behaviour for Ab(c) is called fault model

- $\{I_1\}$  and  $\{I_2\}$  still diagnoses,
- $\{I_1, I_2\}$  no longer a diagnosis, because

 $\mathrm{SD} \cup \mathrm{OBS} \cup {\mathrm{Ab}(I_1), \mathrm{Ab}(I_2)} \vDash \bot$ 

### Conclusions

- Consistency-based diagnosis popular for trouble shooting of equipment and devices
- Extensions: temporal behaviour and continuous behaviour
- Diagnoses may be ranked using probability theory and entropy (General Diagnostic Engine, GDE)

#### Software:

- GDE/ATMS (Palo Alto Research Center): http://www2.parc.com/spl/members/dekleer/
- Leancop (University Darmstad/Potsdam): http://www.leancop.de/
- AILog!