

## Lecture 1. Introduction, syntax and operational semantics of untyped lambda calculus

1. Make exercises 2.5 – 2.10 of *Introduction to Lambda Calculus - selected pages* by Barendregt and Barendsen.

## Lecture 2. Simple type theory, formulas-as-types and proofs-as-terms

See the course notes – notably *Introduction to Type Theory* by Herman Geuvers – and the slides of the lecture.

Do the exercises that are appropriate for your level.

1. Verify in detail (by giving a derivation in  $\lambda \rightarrow$ ) that

$$\lambda x^{\alpha \rightarrow \beta}. \lambda y^{\beta \rightarrow \gamma}. \lambda z^{\alpha}. y(xz) : (\alpha \rightarrow \beta) \rightarrow (\beta \rightarrow \gamma) \rightarrow \alpha \rightarrow \gamma$$

2. (a) Verify in detail (by giving a derivation in  $\lambda \rightarrow$ ) that

$$\lambda x^{\beta \rightarrow \alpha}. \lambda y^{(\beta \rightarrow \alpha) \rightarrow \alpha}. y(\lambda z^{\beta}. xz) : (\beta \rightarrow \alpha) \rightarrow ((\beta \rightarrow \alpha) \rightarrow \alpha) \rightarrow \alpha$$

- (b) “Dress up” the  $\lambda$ -term  $\lambda x. \lambda y. y(\lambda z. xz)$  with type information in such a way that it is of type  $(\beta \rightarrow \gamma) \rightarrow ((\beta \rightarrow \gamma) \rightarrow \alpha) \rightarrow \alpha$
- (c) Give a “simpler” term of type  $(\beta \rightarrow \gamma) \rightarrow ((\beta \rightarrow \gamma) \rightarrow \alpha) \rightarrow \alpha$ .

3. (a) Give the natural deduction (either in Fitch style or in tree form) that corresponds to

$$\lambda x: \gamma \rightarrow \varepsilon. \lambda y: (\gamma \rightarrow \varepsilon) \rightarrow \varepsilon. y(\lambda z: \gamma. yx) : (\gamma \rightarrow \varepsilon) \rightarrow ((\gamma \rightarrow \varepsilon) \rightarrow \varepsilon) \rightarrow \varepsilon$$

- (b) Give another term of the same type

$$(\gamma \rightarrow \varepsilon) \rightarrow ((\gamma \rightarrow \varepsilon) \rightarrow \varepsilon) \rightarrow \varepsilon$$

and the natural deduction (either in Fitch style or in tree form) that it corresponds to.

4. In all of the following cases: give a typing derivation.

- (a) Find a term of type  $(\delta \rightarrow \delta \rightarrow \alpha) \rightarrow (\alpha \rightarrow \beta \rightarrow \gamma) \rightarrow (\delta \rightarrow \beta) \rightarrow \delta \rightarrow \gamma$

- (b) Find two terms of type  $(\delta \rightarrow \delta \rightarrow \alpha) \rightarrow (\gamma \rightarrow \alpha) \rightarrow (\alpha \rightarrow \beta) \rightarrow \delta \rightarrow \gamma \rightarrow \beta$
- (c) Find a term of type  $((\alpha \rightarrow \beta) \rightarrow \alpha) \rightarrow (\alpha \rightarrow \alpha \rightarrow \beta) \rightarrow \alpha$
5. Consider the following term “with holes”  $N$ , where  $A = \alpha \rightarrow \alpha$  and  $\mathbf{I}_1$  and  $\mathbf{I}_2$  and  $\mathbf{I}_3$  and  $\mathbf{I}_4$  are copies of the well-known  $\lambda$ -term  $\mathbf{I} (= \lambda x.x)$ .

$$N \quad := \quad \lambda y:?.(\lambda x:A \rightarrow A. \mathbf{I}_1 (x \mathbf{I}_4 (\mathbf{I}_3 y))) \mathbf{I}_2$$

Fill in the type for ? in  $N$ , give the types for  $\mathbf{I}_1$  and  $\mathbf{I}_2$  and  $\mathbf{I}_3$  and  $\mathbf{I}_4$  and give the type of  $N$  itself in simple type theory ( $\lambda \rightarrow$ ) à la Church. (Note that  $A$  abbreviates  $\alpha \rightarrow \alpha$ .)

### Lecture 3. First order dependent type theory, formulas-as-types and proofs-as-terms

NB.  $\rightarrow$  binds strongest.

1. Give a precise derivation of the following judgment.

$$A : *, P : A \rightarrow *, a : A \vdash (P a) \rightarrow * : \square$$

(Advise: give the derivation in “flag style”, as it was shown in the lecture.)

2. Find a term of the following type and write down the context in which this term is typed.

$$(\Pi x:A. P x \rightarrow Q x) \rightarrow (\Pi x:A. P x) \rightarrow \Pi x:A. Q x$$

Do this by giving a derivation in “flag style”, where you may omit derivations of the well-formedness of types.

3. Find a term of the following type and write down the context in which this term is typed.

$$(\Pi x:A. P x \rightarrow \Pi z:A. R x z) \rightarrow (\Pi x:A. P x) \rightarrow \Pi z:A. R z z.$$

(NB. Read this type in the proper way:  $\rightarrow$  binds stronger than  $\Pi$ !)

4. Give a term  $M$  of type  $\Pi x:A. P(f(f x))$  in the context

$$\begin{aligned} \Gamma \quad := \quad & A : *, P : A \rightarrow *, f : A \rightarrow A, g : A \rightarrow A, \\ & h : \Pi x:A. P(f x) \rightarrow P(g x), k : \Pi x, y:A. (P x \rightarrow P y) \rightarrow P(f x). \end{aligned}$$

Also give a derivation of  $\Gamma \vdash M : \Pi x:A. P(f(f x))$  in ‘short form’, so you don’t have to show the well-formedness of the types.

5. Find a term of the following type and write down the context in which this term is typed.

$$(\Pi x:A. P x \rightarrow Q) \rightarrow (\Pi x:A. P x) \rightarrow Q$$

What is special about your context? Explain how your context explicitly ensures a property for the type  $A$ .

6. Find a term from the given hypotheses of the following type and write down the context in which this term is typed.

$$\begin{aligned} & \forall x. (P(x) \rightarrow R(x, f(x))), \\ & \forall x, y. (R(x, y) \rightarrow R(y, x)), \\ & \forall x, y. (R(x, y) \rightarrow R(f(y), x)) \quad \vdash \quad \forall x. (P(x) \rightarrow R(f(x), f(x))) \end{aligned}$$