Lambda-Calculus and Type Theory ISR 2024

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Exercises Day 4

lecture 9. Church-Rosser property

All exercises are about the Church-Rosser proof that is on the slides (due to Takahashi) and that we have presented at the lecture. We recall the definition of \Rightarrow using derivation rules:

$$\frac{1}{x \Rightarrow x} \text{ (var)} \qquad \frac{M \Rightarrow M'}{\lambda x.M \Rightarrow \lambda x.M'} (\lambda)$$

$$\frac{M \Rightarrow M' \qquad N \Rightarrow N'}{M \ N \Rightarrow M' \ N'} \text{ (app)} \qquad \frac{M \Rightarrow M' \qquad N \Rightarrow N'}{(\lambda x. M) \ N \Rightarrow M'[N'/x]} (\beta)$$

- 1. Consider the term $M = (\lambda x y.x x(x y))(\mathbf{I} \mathbf{I})$
 - (a) Give the reduction graph of M. (You may abbreviate **II** to J and $\lambda x y.x x(x y)$ to P.)
 - (b) Compute M^* and $(M^*)^*$.
 - (c) Prove that $M \Rightarrow M^*$ and $M^* \Rightarrow (M^*)^*$ by giving a derivation.
- 2. In the definition of \Rightarrow , we change clause (β) into

$$\frac{M \Rightarrow \lambda x.P \qquad N \Rightarrow N'}{M \, N \Rightarrow P[N'/x]}$$

- (a) Give the definition of $(-)^*$ that goes with this adapted definition of \Rightarrow .
- (b) Prove again (with these adapted definitions) that $M \Rightarrow N$ implies $N \Rightarrow M^*$, by doing the inductive step for case (β) .
- 3. The η -reduction rule is: $\lambda x.M \ x \to_{\eta} M$, if $x \notin FV(M)$. In order to prove CR for $\beta \eta$ we add a clause for η -redexes to the definition of \Rightarrow :

$$\frac{M \Rightarrow M'}{\lambda x.M \, x \Rightarrow M'} \, x \notin \mathrm{FV}(M)$$

- (a) Show that now $(\lambda yx.yx)\mathbf{I} \Rightarrow \mathbf{I}$, and show that in the original definition, this is not the case.
- (b) Define $(-)^*$ for this extension to η