Lambda-Calculus and Type Theory ISR 2024 Obergurgl, Austria Herman Geuvers & Niels van der Weide Radboud University Nijmegen NL Exercises Day 6

Lecture 12. Principal types: a functional programmers' view on type theory

- 1. (a) Determine the most general unifier of $(\alpha \rightarrow \beta) \rightarrow \gamma$ and $\alpha \rightarrow \beta \rightarrow \gamma$.
 - (b) Determine the most general unifier of $(\alpha \rightarrow \beta) \rightarrow \gamma$ and $\gamma \rightarrow \alpha \rightarrow \beta$
- 2. Compute the principal type of $\mathbf{S} := \lambda x \cdot \lambda y \cdot \lambda z \cdot x z(y z)$.
- 3. Consider the following two terms
 - $\lambda x.x (\lambda y.y (\lambda z.x))$
 - $\lambda x.x (\lambda y.x (\lambda z.z))$

For each of these terms, compute its principal type, if it exists. (Give the end result and show your computation; if the term has no principal type, show how your computation yields 'fail'.)

- 4. For each of the following two terms, compute its principal type, if it exists.
 - $\lambda x.(\lambda y.x(xy))(\lambda u v.u)$
 - $\lambda y.(\lambda x.x(xy))(\lambda u v.u)$

Give the end result and show your computation; if the term has no principal type, show how your computation yields 'fail'.