

Lambda-Calculus and Type Theory  
ISR 2024  
Oberurgl, Austria  
Herman Geuvers & Niels van der Weide  
Radboud University Nijmegen NL

**Test**

The maximum number of points per exercise is indicated in the margin. Maximum 100 points in total.  
**NB** Typing derivations may be given in “flag style” or in “sequent style”.  
**You can do exercises 7 and 8 either in Coq or on paper. In case you do these in Coq, use the test.v file on the website and send the completed file to herman@cs.ru.nl and nweide@cs.ru.nl**

---

- (10) 1. Give a typing derivation that shows that the following term is typable in simple type theory  $(\lambda \rightarrow)$  à la Church.

$$\lambda x : \alpha. \lambda y : ((\alpha \rightarrow \alpha) \rightarrow \alpha) \rightarrow \alpha. y (\lambda z : \alpha \rightarrow \alpha. z x)$$

- (10) 2. Give a term  $M$  in simple type theory  $(\lambda \rightarrow)$  à la Church with type

$$(\gamma \rightarrow \beta) \rightarrow ((\alpha \rightarrow \beta) \rightarrow \alpha) \rightarrow \gamma \rightarrow \alpha.$$

Give a typing derivation that gives the type of  $M$ .

3. Consider the following term in the simply typed  $\lambda$ -calculus.

$$N \quad := \quad (\lambda y : \alpha \rightarrow \alpha. \lambda x : \alpha. y (y x)) \quad (\mathbf{II})$$

- (8) (a) Compute  $m(N)$ , where  $m(-)$  is the measure of a term, as used in the proof of weak normalization (WN) for  $\lambda \rightarrow$ . Indicate which redex will be contracted according to the strategy described in the WN proof.
- (3) (b) Now consider  $N$  as a term in the untyped  $\lambda$ -calculus. Calculate  $N^*$ . (You can just give  $N^*$ , you don't have to show a computation.)
- (6) (c) Give all terms  $M$  for which we have  $N \Rightarrow M$ . (You don't have to give the derivations of  $N \Rightarrow M$ .)

- (10) 4. Show that, in  $\lambda 2$  à la Church,

$$\lambda x : \perp. x ((\top \rightarrow \top \rightarrow \top) \rightarrow \top) (\lambda y : \top. y \top) \quad : \quad \perp \rightarrow \top,$$

where  $\top := \forall \alpha : *. \alpha \rightarrow \alpha$  and  $\perp := \forall \alpha : *. \alpha$ . Give the typing derivation of your result.

**Continue on the other side**

- (12) 5. In  $\lambda 2$  à la Church, we have the well-known types of natural numbers and of lists over natural numbers:

$$\begin{aligned}\mathbb{N} &:= \forall \alpha : * . \alpha \rightarrow (\alpha \rightarrow \alpha) \rightarrow \alpha \\ \mathbb{L}_{\mathbb{N}} &:= \forall \alpha : * . \alpha \rightarrow (\mathbb{N} \rightarrow \alpha \rightarrow \alpha) \rightarrow \alpha.\end{aligned}$$

You may assume that the function  $\mathbf{plus} : \mathbb{N} \rightarrow \mathbb{N} \rightarrow \mathbb{N}$  (that adds two natural numbers) is known. Define the function  $\mathbb{L}\mathbf{plus} : \mathbb{L}_{\mathbb{N}} \rightarrow \mathbb{N}$  that adds the numbers in a list.

- (10) 6. In the system  $\lambda P$ , give a term of type  $\Pi x:A.R(f a)(f(f x))$  in the context

$$\begin{aligned}\Gamma &:= A : *, a : A, R : A \rightarrow A \rightarrow *, Q : A \rightarrow A \rightarrow *, f : A \rightarrow A, \\ p &: \Pi x:A.Q a(f x), \quad q : \Pi x, y:A.Q x y \rightarrow R x y, \\ r &: \Pi x, y:A.R x y \rightarrow R(f x)(f y).\end{aligned}$$

That is, solve the following type inhabitation problem in  $\lambda P$ :

$$\Gamma \vdash ? : \Pi x:A.R(f a)(f(f x))$$

Also give a typing derivation in short version (where you don't have to establish the well-formedness of the types themselves).

- (15) 7. [You can also do this exercise in Coq] In the system  $\lambda P$ , give a term of type

$$\Pi x, y:A.R x y \rightarrow R(f x) x$$

in the context

$$\begin{aligned}\Gamma &:= A : *, R : A \rightarrow A \rightarrow *, f : A \rightarrow A \\ h &: \Pi x, y:A.R x y \rightarrow R(f y) x \\ t &: \Pi x, y, z:A.R x y \rightarrow R(f y) z \rightarrow R x z\end{aligned}$$

You don't have to give a typing derivation!

8. [You can also do this exercise in Coq] In Lecture 7 we discussed Higher order logic and the Calculus of Constructions. We defined, given  $A : *, P : A \rightarrow *, f : A \rightarrow A$ , the *smallest subset of  $A$  containing  $P$  and closed under  $f$*  as follows.

$$S := \lambda y : A. \forall Q : A \rightarrow *. (P \subseteq Q) \rightarrow (\forall x : A. Q x \rightarrow Q(f x)) \rightarrow Q y$$

where  $P \subseteq Q := \forall x : A. P x \rightarrow Q x$ .

In the lecture we proved that  $S$  is closed under  $f$ .

- (8) (a) Prove that  $S$  contains  $P$  (that is, give a term of type  $P \subseteq S$ ).
- (8) (b) Prove that  $S$  is indeed the smallest such (that is give a term that proves that for all  $Q$  that contain  $P$  and are closed under  $f$ , we have  $S \subseteq Q$ ).

You don't have to give a typing derivation!

**END**

---