

# Coalgebra lecture 12, exercise 5

December 3, 2018

Let  $(X, \rightarrow)$  be a labelled transition system over  $A$ . Let  $A^\infty = A^\omega \cup A^*$  be the set of streams and (finite) words over  $A$ . The empty word is denoted by  $\varepsilon$ . Consider the following rules, involving a relation  $\downarrow \subseteq X \times A^\infty$ .

$$\frac{}{x \downarrow \varepsilon} \qquad \frac{x \xrightarrow{a} y \quad y \downarrow w}{x \downarrow aw} \qquad (1)$$

(for all  $a \in A$ ,  $w \in A^\infty$ ).

Let  $\text{Rel}_{X, A^\infty}$  be the set of relations of the form  $R \subseteq X \times A^\infty$ , partially ordered by subset inclusion  $\subseteq$ . This partial order is a complete lattice.

Given  $x \in X$  and  $w \in A^\infty$ , we say  $w$  is a *trace* of  $x$  if there is a path from  $x$  labelled by  $w$ , that is, a path  $x_1 \xrightarrow{a_1} x_2 \xrightarrow{a_2} x_3 \xrightarrow{a_3} \dots$  such that  $x_1 = x$  and  $w = a_1 a_2 a_3 \dots$ . If  $w \in A^*$  then we call this a *finite* trace, if  $w \in A^\omega$  we call it an *infinite* trace.

- (a) Describe the least upper bound  $\bigvee S$  and greatest lower bound  $\bigwedge S$  of an arbitrary set  $S \subseteq \text{Rel}_{X, A^\infty}$ , and give the top and bottom elements of the lattice (you don't have to give a proof).

*Solution.* The least upper bound of a set  $S \subseteq \text{Rel}_{X, A^\infty}$  is given by union:  $\bigvee S = \bigcup_{R \in S} R$ , the greatest lower bound by intersection  $\bigwedge S = \bigcap_{R \in S} R$ , the top element by  $X \times A^\infty$  and the bottom element by  $\emptyset$ .

- (b) Formulate the rules (1) in terms of a function  $b: \text{Rel}_{X, A^\infty} \rightarrow \text{Rel}_{X, A^\infty}$ . Show that your function is monotone.

*Solution.*

$$b(R) = \{(x, w) \mid w = \varepsilon \text{ or } \exists a \in A, v \in A^\infty, y \in X. w = av, x \xrightarrow{a} y \text{ and } (v, y) \in R\}.$$

For monotonicity, suppose  $R \subseteq S$ ; we should prove that  $b(R) \subseteq b(S)$ . Let  $(x, w) \in b(R)$ . If  $w = \varepsilon$  then  $(x, w) \in b(S)$  is immediate from the definition; if  $w = av$  then  $x \xrightarrow{a} y$  for some  $y$  with  $(y, v) \in R$ . But  $R \subseteq S$ , so  $(y, v) \in S$ , and it follows that  $(x, w) \in b(S)$ .

- (c) What is a pre-fixed point of  $b$ ? And what is the least fixed point? Give a concrete description, in terms of the transition system and elements of  $A^\infty$ .

*Solution.* A relation  $R$  is a pre-fixed point of  $b$  if  $b(R) \subseteq R$ . Concretely, this means

- $(x, \varepsilon) \in R$  for all  $x \in X$ , and
- if  $x \xrightarrow{a} y$  and  $(w, y) \in R$  then  $(x, aw) \in R$ .

The least fixed point is also the least pre-fixed point: the least relation satisfying the above conditions. It is given by the relation  $\{(x, w) \mid w \text{ is a finite trace of } x\}$ .

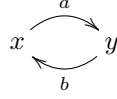
- (d) What is a post-fixed point of  $b$ ? And what is the greatest fixed point? Give a concrete description, in terms of the transition system and elements of  $A^\infty$ .

*Solution.* A relation  $R$  is a post-fixed point of  $b$  if  $R \subseteq b(R)$ . This means:  $\forall (x, w) \in R$ ,

- $w = \varepsilon$ , or
- $w = av$ ,  $x \xrightarrow{a} y$  and  $(y, v) \in R$ .

The greatest fixed point is the greatest post-fixed point, and is given by the relation  $\{(x, w) \mid w \text{ is a (finite or infinite) trace of } x\}$ .

- (e) Use your answer to one of the previous questions to show that, in the transition system below, every finite trace of  $x$  is a prefix of the stream  $ababab \dots$



*Solution.* As explained above, the finite traces are given by  $\text{lfp}(b)$ , the least fixed point of  $b$ . The induction proof principle (following from Knaster-Tarski) states that

$$\frac{b(R) \subseteq R}{\text{lfp}(b) \subseteq R} \quad (2)$$

for every  $R \in \text{Rel}_{X, A^\omega}$ . We choose

$$R = \{(x, w) \mid w \text{ is a prefix of } (ab)^\omega\} \cup \{(y, w) \mid w \text{ is a prefix of } (ba)^\omega\}.$$

It is easy to check that  $R$  is indeed a pre-fixed point of  $b$ , hence  $\text{lfp}(b) \subseteq R$ ; so whenever  $w$  is a trace of  $x$ , we have  $(x, w) \in R$ , which means  $w$  is a prefix of  $(ab)^\omega$ .

- (f) Use your answer to one of the previous questions to show that, in the transition system above, the stream  $ababab \dots$  is an infinite trace of  $x$ .

*Solution.* The relation of finite and infinite traces is given by  $\text{gfp}(b)$ , the greatest fixed point of  $b$ . Thus, it suffices to show that  $(x, (ab)^\omega) \in \text{gfp}(b)$ . To do so, we use the coinductive proof principle:

$$\frac{R \subseteq b(R)}{R \subseteq \text{gfp}(b)} \quad (3)$$

and choose

$$R = \{(x, (ab)^\omega), (y, (ba)^\omega)\}.$$

It is easy to check that  $R$  is a post-fixed point of  $b$ , so that  $R \subseteq \text{gfp}(b)$ , which in particular means  $(x, (ab)^\omega) \in \text{gfp}(b)$ .

- (g) Suppose that our transition system is finitely branching, meaning that for each  $x \in X$ , the set  $\{y \mid x \xrightarrow{a} y \text{ for some } a\}$  is finite. Consider the relation  $\parallel \subseteq X \times A^\infty$ , given by:  $x \parallel w$  iff there are infinitely many prefixes that are finite traces of  $x$ . Prove that for all  $x \in X$  and  $w \in A^\infty$ : if  $x \parallel w$ , then  $w$  is an infinite trace of  $x$ .

*Solution.* We are going to prove that  $\parallel$  is a post-fixed point of  $b$ . Again by coinduction (3), we then obtain that  $x \parallel w$  implies  $(x, w) \in \text{gfp}(b)$ , meaning that  $w$  is a trace of  $x$ .

Suppose  $x \parallel w$ . We need to prove that  $(x, w) \in b(\parallel)$ . First observe that  $w \in A^\omega$ , since it has infinitely many prefixes. Hence  $w = av$  for some  $a \in A$ . Now, by definition of traces, for

every trace of  $x$  of the form  $au$  there exists a state  $y$  such that  $x \xrightarrow{a} y$ , and  $u$  is a trace of  $y$ . In particular, there are infinitely many prefixes of  $w = av$  which are traces of  $x$ , so there are infinitely many prefixes of  $v$  which are traces of  $y$  for some state  $y$  with  $x \xrightarrow{a} y$ . But there are only *finitely* many  $y$  such that  $x \xrightarrow{a} y$ , by the assumption that the LTS is finitely branching. Hence, there must be a state  $y$  such that  $x \xrightarrow{a} y$  and infinitely many prefixes of  $v$  are a trace of  $y$ . Thus  $y \parallel v$ , and  $(x, y) \in b(\parallel)$ .